# Sign Patterns, Nonsingularity, and the Solvability of $\boldsymbol{A x}=\boldsymbol{b}$ 

Charles R. Johnson*<br>Institute for Physical Science and Technology<br>and Department of Economics<br>University of Maryland<br>College Park, Maryland 20742<br>Frank Uhlig<br>Institut fur Geometrie und Praktische Mathematik<br>RWTH Aachen<br>51 Aachen, West Germany

and

Dan Warner
Mathematical Sciences Department
Clemson University
Clemson, South Carolina 29631

Submitted by George P. Barker


#### Abstract

We investigate conditions on the sign pattern class of the ( $n-1$ )st compound of a real $n$-by- $n$ matrix $A$ such that the solvability of $A x=b^{(i)}$ for $i=1, \ldots, k, k<n$, with specific $b^{(i)}$, insures the nonsingularity of $A$. The number and choice of right-hand sides $b^{(i)}$ sufficient for the task depends only on the sign-pattern class of the $(n-1)$ st compound of $A$. The result for $k=1$ generalizes a known fact about totally nonnegative matrices and an observation about $M$-matrices, thus providing another unifying result for these two classes of matrices.


[^0]
## INTRODUCTION

Let $A=\left(a_{i j}\right)$ be an $n$-by- $n$ real matrix. We raise the following question. What information about $A$ permits the deduction that $A$ is nonsingular from knowing that

$$
\begin{equation*}
A x=b, \quad b \neq 0 \tag{1}
\end{equation*}
$$

can be solved for a specified set of $k<n$ right-hand sides $b$ ? This question is motivated by the intriguing, but not too well-known, fact that if $A$ is totally nonnegative (the determinant of each square submatrix is nonnegative), then the existence of an $x$ such that the components of $b=A x$ alternate in sign implies that $A$ is nonsingular [2]. We refer to this as the "totally nonnegative result," but actually somewhat more can be said. In applications, the totally nonnegative hypothesis is often assured by the structure of a given problem, so that the result takes on added significance.

Our interest is a theoretical one. We try to generalize, and in the process re-prove, the totally nonnegative result by relating sign patterns and the solvability of (1) to the nonsingularity of $A$. [The nonsingularity of $A$, of course, may, in principle, be determined in the same degree of complexity as a solution to (1).] In this spirit, our focus is on information about the minors, most specifically the signs of the $(n-1)$-by- $(n-1)$ minors of $A$. It is clear that, if highly specific information is known, then little more may need to be done in order to verify nonsingularity for $A$. For example, if it is known that any particular $(n-1)$-by- $(n-1)$ submatrix has nonzero determinant, say that achieved by deletion of row $i$ and column $j$, then $A$ is nonsingular if and only if (1) is solvable for $b$ equal to the $i$ th unit vector (Cramer's rule). However, such specific information may not be available, and our sign-pattern generalizations of the totally nonnegative result, which include the above observation as a special case, do not generally assume that any particular minor is nonzero. They do, however, usually assume that some $(n-1)$-by- $(n-1)$ minor is nonzero, for if all were zero, then $A$ would be singular.

## NOTATION AND CONCEPTS

For two index sets $I, J \subseteq\{1,2, \ldots, n\}$, we denote by $A(I, J)$ the submatrix of $A$ resulting from deletion of the rows indicated by $I$ and the columns indicated by $J$. When $I=\{i\}$ and $J=\{j\}$, each contains exactly one index, and we abbreviate $A(\{i\},\{i\})$ to $A(i, j)$. A $k$-by- $k$ minor of $A$ is a scalar $\operatorname{det} A(I, J)$, where $I=\left\{i_{1}, \ldots, i_{n-k}\right\}$ and $J=\left\{i_{1}, \ldots, i_{n-k}\right\}$, and a minor is said
to be principal if it is the determinant of a principal submatrix, i.e. if $I=J$. We let $C_{k}(A)$ denote the $k t h$ compound matrix of $A$ : the $\binom{n}{k}$-by- $\binom{n}{k}$ matrix of $k$-by- $k$ minors of $A$, ordered lexicographically. If $F$ is the nonsingular diagonal matrix, $F \equiv \operatorname{diag}\left((-1)^{i}\right)$, whose diagonal entries alternate between -1 and +1 , then the adjoint matrix of $A$ is defined by

$$
\begin{equation*}
\operatorname{adj} A \equiv F C_{n-1}(A)^{T} F \tag{2}
\end{equation*}
$$

Of course, $A^{-1}=(\operatorname{det} A)^{-1} \operatorname{adj} A$ if $A$ is nonsingular, and otherwise, $\operatorname{adj} A$ [and $C_{n-1}(A)$ ] is of $\operatorname{rank} 1$ (if $\operatorname{rank} A=n-1$ ) or $\operatorname{rank} 0$ (if $\operatorname{rank} A \leqslant n-2$ ). We note at this point that, although we state our results in terms of the compound $C_{n-1}(A)$, it will be clear that there are equivalent statements in terms of adj A. Finally, we call a real entried vector or matrix uniformly signed if all entries are nonnegative or all entries are nonpositive. A particular uniformly signed vector is the vector $e$, each of whose entries is equal to 1 . We define the alternating-sign vector to be $f \equiv \mathrm{Fe}$.

We utilize two notions of a sign-pattern matrix and associated sign-pattern class. A weak-sign-pattern matrix $P$ may have entries of four possible types: " + " denotes a nonnegative entry, " -" a nonpositive entry, " 0 " allows only the entry zero, and "*" indicates unrestricted entries. Such a sign-pattern matrix $P$ defines a class of real matrices $\mathscr{P}$ in a natural way, except that we make the additional requirement that for $A$ to be in $\mathscr{P}$, not all entries of $A$ corresponding to + 's and -'s in $P$ may be zero. Note then that $\mathscr{P}$ is empty if $P$ has no entries equal to + or - . For example, if

$$
P=\left[\begin{array}{cc}
+ & 0 \\
* & -
\end{array}\right]
$$

then

$$
\left[\begin{array}{rr}
1 & 0 \\
-2 & -5
\end{array}\right],\left[\begin{array}{rr}
0 & 0 \\
2 & -1
\end{array}\right] \in \mathscr{\varphi}
$$

while

$$
\left[\begin{array}{ll}
0 & 0 \\
3 & 0
\end{array}\right] \text { and }\left[\begin{array}{rr}
-1 & 0 \\
2 & -1
\end{array}\right]
$$

are not. Our results will be based upon sign-pattern information about $C_{n-1}(A)$ and will thus apply to classes of matrices.

Our second notion of a sign pattern is that of a strong-sign-pattern matrix, which is, in a certain sense, dual to the weak-sign-pattern notion mentioned above. The sign-pattern matrix $Q=P^{d}$ has " + " entries and " -" entries wherever $P$ does, but these two symbols now require the strong interpretation of positive or negative entries, respectively. The remaining entries of $Q$ are determined by replacing the "*" entries of $P$ with 0 's and the " 0 " entries with *'s. These two symbols are interpreted as before, and, again, $Q$ naturally determines a class of real matrices, which we denote by 2 . For example, if

$$
P=\left[\begin{array}{ll}
* & + \\
0 & -
\end{array}\right]
$$

then

$$
Q=P^{d}=\left[\begin{array}{ll}
0 & + \\
* & -
\end{array}\right]
$$

is the strong-sign-pattern matrix associated with $P$, and

$$
\left[\begin{array}{rr}
0 & 3 \\
-1 & -2
\end{array}\right] \in \mathscr{2}, \quad \text { while } \quad\left[\begin{array}{rr}
0 & 2 \\
-1 & 0
\end{array}\right] \notin \mathscr{Q} .
$$

Notationally, if $Q=P^{d}$, we also write $\mathcal{Q}=\mathscr{P}^{d}$. By construction, the strong sign-pattern class $\mathcal{Q}=\mathscr{P}^{d}$ is such that the trace of each $A B^{T}, A \in \mathscr{P}, B \in \mathcal{Q}$, is positive, and this is the sense of duality mentioned above. The requirements placed upon 2 are minimally necessary to insure this.

By the rank of a strong-sign-pattern matrix $Q$, we mean the minimum of the ranks assumed by all matrices in the class $\mathcal{Q}$, and we shall say that the rank of a weak sign-pattern matrix $P$ is $k$ if $k=\operatorname{rank} P^{d}$. There will always be matrices $B \in \mathscr{Q}$ such that $\operatorname{rank} B=\operatorname{rank} Q$, but the rank of any matrix in $\mathscr{Q}$ serves as an upper bound for rank $Q$. In particular, the special matrix $C(Q)$, in which +'s in $Q$ are replaced by l's, -'s by - l's, and 0 's or *'s by 0 's, provides an upper bound, which makes it clear that

$$
\left[\begin{array}{ccc}
+ & \cdots & + \\
\vdots & & \vdots \\
+ & \cdots & +
\end{array}\right] \text { and }\left[\begin{array}{ccccc}
+ & - & + & - & \cdots \\
- & & & & \\
\vdots & & & &
\end{array}\right]
$$

are rank-1 sign patterns. It is an interesting question how to determine the rank of a sign pattern. Note that multiplication of any strong sign pattern by a
diagonal strong-sign-pattern matrix (a signature matrix) is well defined, and such a transformation (no zero diagonal entries) on the right or left, and left or right multiplication by a permutation matrix, as well as transposition, do not change the rank of a strong sign pattern. Another interesting question is to determine the equivalence classes of strong-sign-pattern matrices under these transformations collectively. The rank of each sign pattern in such an equivalence class is the same.

## PRIMARY RESULT

An observation which underlies our results is the following very simple fact.

Lemma. If, for un n-by-n mutrix A, there exists an $x$ such that (1) is satisfied for a given $b$, then

$$
\begin{equation*}
x_{i} \operatorname{det} A=\sum_{i=1}^{n}(-1)^{i+i} b_{i} \operatorname{det} A(i, j) \tag{3}
\end{equation*}
$$

for each $;=1, \ldots, n$.

Proof. This is essentially Cramer's rule before dividing by det $A$. In general, $(\operatorname{adj} A) A=(\operatorname{det} A) I$, so that $(\operatorname{adj} A) b=(\operatorname{adj} A) A x=(\operatorname{det} A) x$, of which (3) is a component-by-component version based upon (2).

Another vectorial version of (3) is

$$
\begin{equation*}
(\operatorname{det} A)(F x)^{T}=b^{T} F C_{n-1}(A) \tag{4}
\end{equation*}
$$

The utility of (3) occurs when the right-hand side can be shown to be nonzero for some $j$; for then the existence of a solution $x$ of (3) implies $\operatorname{det} A \neq 0$. Alternatively stated, if (1) is solvable and $A$ is singular, then $b$ must be orthogonal to every column of $F C_{n-1}(A)$. [Conversely, if $b^{T} F C_{n-1}(A)=0$, then $A$ is singular.]

Our principal observation is the

Theorem. Let $\mathscr{P}$ be a rank $k \geqslant 1$ weak-sign-pattern class. Then there exist $k$ vectors

$$
b^{(1)}, b^{(2)}, \ldots, b^{(k)}
$$

which depend only upon the class $\mathscr{P}$, such that if $C_{n-1}(A) \in \mathscr{P}$, then $A$ is nonsingular if and only if $(1)$ is solvable for each $b=b^{(i)}, i=1, \ldots, k$.

Proof. Choose $B \in \mathcal{Q}=\mathscr{P}^{d}$, the strong-sign-pattern class, such that rank $B=k$, and let $b^{(1)}, \ldots, b^{(k)}$ be $k$ linearly independent columns of FBF. We show that these $b^{(i)}$ satisfy the requirements of the theorem. Suppose, as we may by the definition of $\mathscr{P}$, that column $j$ of $C_{n-1}(A)$ contains a nonzero entry in a position corresponding to a + or - in the matrix $P$ which defines $G$, and consider column $j$ of $F B F$, which we call $\hat{b}$. If (1) is solvable for each right-hand side $b^{(i)}, i=1, \ldots, k$, then (1) is solvable for $b=\hat{b}$, since it is in the span of the $b^{(i)}$. By virtue of the construction of 2 , if the $i$ th entry of column $j$ of $C_{n-1}(A)$ is positive, the corresponding entry of $\hat{b}$ has the sign of $(-1)^{i+j}$ or is 0 , and, if the $i$ th entry of column $j$ of $C_{n-1}(A)$ is negative, the corresponding entry of $\hat{b}$ has the sign of $(-1)^{i+j+1}$ or is 0 . Furthermore, there is at least one nonzero entry in column $\boldsymbol{j}$ of $C_{n-1}(A)$ which corresponds to a + or - in $P$, and any such nonzero entry must correspond to a nonzero entry in $\hat{b}$; therefore, since $\mathscr{2}$ is a strong-sign-pattern class, the right-hand side of (3) is positive for this $j$ and $b=\hat{b}$, and it follows that det $A \neq 0$, since (1) is solvable for $b=\hat{b}$. Thus the solvability of (1) for each $b=b^{(i)}, i=1, \ldots, k$, implies that $A$ is nonsingular. If $A$ is nonsingular, then (1) is, of course, solvable for any right-hand side $b^{(i)}$ and the proof is complete.

## IMPLICATIONS

The case $k=1$ of the theorem is of note because the fewest right-hand sides are required. Moreover, rank-1 patterns are easily recognized.

Remark. Let $J$ be the strong sign pattern

$$
J \equiv\left[\begin{array}{ccc}
+ & \cdots & + \\
\vdots & & \vdots \\
+ & \cdots & +
\end{array}\right]
$$

It is clear that rank $J=1$. Let $D_{1}$ and $D_{2}$ be two diagonal sign-pattern matrices with the restriction that each has at least $\mathbf{a}+$ or - on the diagonal. It is furthermore clear that the sign-pattern product

$$
\begin{equation*}
Q=D_{1} J D_{2} \tag{5}
\end{equation*}
$$

has rank 1 , and it is easy to check that any,,+- 0 rank- 1 class is determined by a sign-pattern matrix of the form (5). Therefore, a strong-sign-pattern matrix $Q$ has rank 1 if and only if the particular matrix $C(Q)$ is of rank 1 . Thus, it is particularly easy to determine if a pattern $Q$ has rank 1 . The rank of $C(Q)$ may be calculated directly, or, since 0 -entries of rank-1 patterns can occur only in 0 -rows or 0 -columns, first any 0 -rows or columns must be sorted out of $Q$, and then, when the remaining rows and columns are sign scaled so that their initial entries are,$+ Q$ is of rank 1 if and only if the result is uniformly signed ( + ).

Example. Although $C(Q)$ suffices to identify rank-l sign patterns $Q$, it does not in general determine the rank of a sign pattern. For example, if

$$
Q=\left[\begin{array}{lll}
- & - & + \\
+ & + & + \\
- & + & -
\end{array}\right]
$$

then $\operatorname{rank} C(Q)=3$, but

$$
\operatorname{rank} Q=\operatorname{rank}\left[\begin{array}{rrr}
-1 & -4 & 2 \\
4 & 1 & 1 \\
-1 & 1 & -1
\end{array}\right]=2
$$

In fact, the rank of any 3 -by- $3+,-$ sign pattern is at most 2 .
We may record the case $k=1$ of the theorem as

Corollary 1, Let $\mathscr{P}$ be a rank-1 weak-sign-pattern class, and let be the first nonzero column of $F C\left(P^{d}\right)$. If $C_{n-1}(A) \in \mathscr{P}$, then $A$ is nonsingular if and only if (1) is solvable for $b$.

Two special cases of Corollary 1 are the totally nonnegative matrices of rank at least $n-1\left(C_{n-1}(A) \in \mathscr{9}\right.$ for $\left.P=J\right)$ and the (possibly singular) $M$-matrices of rank at least $n-1\left(C_{n-1}(A) \in \mathscr{P}\right.$ for $\left.P=F J F\right)$. Thus, Corollary 1 may be viewed as a unifying result for totally nonnegative matrices and $M$-matrices.

Corollary 2. Suppose that $A$ is an n-by-n real matrix such that $C_{n-1}(A)$ is uniformly signed and $C_{n-1}(A) \neq 0$. If (1) is solvable for $b=f$, then $A$ is nonsingular.

Corollary 3. Suppose that $A$ is an n-by-n real matrix such that $F C_{n-1}(A) F$ is uniformly signed and $C_{n-1}(A) \neq 0$. If (1) is solvable for $b=e$, then $A$ is nonsingular.

Actually, when Corollaries 2 and 3 are specialized further to the cases of totally nonnegative and (possibly singular) $M$-matrices respectively, the assumption $C_{n-1}(A) \neq 0$ (rank at least $n-1$ ) may be omitted. This is because the defining property of each class is inherited under the extraction of principal submatrices and because the vectors $e$ and $f$ retain their form when components are deleted serially, which permits an argument on a principal submatrix in case the rank is smaller than $n-1$. Recall that an $M$-matrix (possibly singular) is one in the weak-sign-pattern class determined by

$$
\mathrm{Z}=\left[\begin{array}{ccccc}
+ & & & & \\
& + & & - & \\
& & \ddots & & \\
& - & & & +
\end{array}\right]
$$

all of whose eigenvalues have nonnegative real parts. An M-matrix is called a nonsingular $M$-matrix if the minimum of these real parts is positive.

Conollany 4. If a real n-by-n matrix $A$ is totally nonnegative, then $A$ is nonsingular if and only if $(1)$ is solvable for $b=f$.

Proof. We proceed by induction. The assertion is clear for $n=1$. If $A x=f$ and $A$ is singular, then there is a $y \neq 0$ such that $A(x+\alpha y)=f$ for all real $\alpha$. Choose $\alpha$ so that $x_{i}+\alpha y_{i}=0$ for some $i$, and define $\tilde{x}$ to be the ( $n-1$ )-vector equal to $x+\alpha y$ with the $i$ th component deleted. It follows that $A(n, i) \tilde{x}=f \in R^{n-1}$. By the induction hypothesis, $A(n, i)$ is nonsingular, $\operatorname{det} A(n, i) \neq 0$, and Corollary 2 applies to $A$. This means that $A$ was actually nonsingular.

Corollary 5. An n-by-n M-matrix A is nonsingular if and only if (1) is solvable for $b=e$.

Proof. This is similar to Corollary 4 except that $A(i, i)$ is used, along with the fact that any component deleted from $e$ leaves a vector of the same form.

Remark. It is clear that all the results mentioned here actually depend only on the sign patterns of the specified right-hand-side vectors (and not on the magnitudes of their components).

## REFERENCES

1 C. R. Johnson, F. T. Leighton, and H. A. Robinson, Sign patterns of inverse positive matrices, Linear Algebra Appl. 24:75-83 (1979).
2 S. Karlin, Total Positivity, Vol. 1, Stanford U.P., Stanford, Calif., 1968, Chapter 5.

Received 13 January 1981; revised 10 October 1981


[^0]:    *The work of this author supported in part by National Science Foundation Grant MCS $80-01611$

